

5.8 all_equal_valley

	DESCRIPTION	LINKS	AUTOMATON
Origin	Derived from valley and all_equal .		
Constraint	<code>all_equal_valley(VARIABLES)</code>		
Argument	VARIABLES : <code>collection(var-dvar)</code>		
Restrictions	$ VARIABLES > 0$ <code>required(VARIABLES, var)</code>		
Purpose	<p>A variable V_k ($1 < k < m$) of the sequence of variables $VARIABLES = V_1, \dots, V_m$ is a <i>valley</i> if and only if there exists an i ($1 < i \leq k$) such that $V_{i-1} > V_i$ and $V_i = V_{i+1} = \dots = V_k$ and $V_k < V_{k+1}$.</p> <p>Enforce all the valleys of the sequence $VARIABLES$ to be assigned the same value, i.e. to be located at the same altitude.</p>		
Example	$((1, 5, 5, 4, 2, 2, 6, 2, 7))$		

The `all_equal_valley` constraint holds since the two valleys, in bold, of the sequence 1 5 5 4 2 2 6 2 7 are located at the same altitude 2. Figure 5.14 depicts the solution associated with the example.

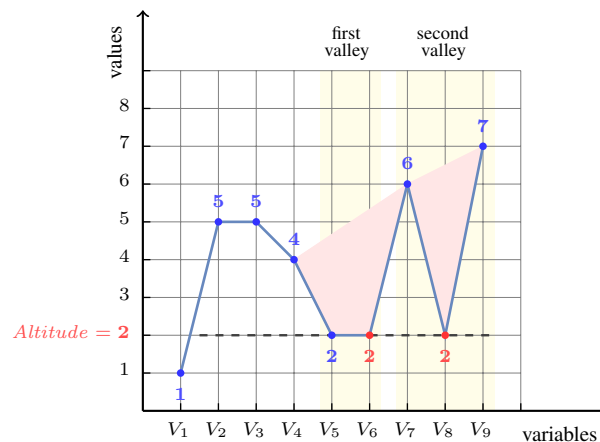


Figure 5.14: Illustration of the **Example** slot: a sequence of nine variables $V_1, V_2, V_3, V_4, V_5, V_6, V_7, V_8, V_9$ respectively fixed to values 1, 5, 5, 4, 2, 2, 6, 2, 7 and its corresponding two valleys, in red, both located at altitude 2

Note that the `all_equal_valley` constraint does not enforce that the minimum value of the sequence $VARIABLES$ corresponds to the altitude of its valleys since, as shown by the

example, the sequence can start with an increasing subsequence that starts below the altitude of its valleys. It also does not enforce that the sequence `VARIABLES` contains at least one valley.

Typical

```
|VARIABLES| ≥ 5
range(VARIABLES.var) > 1
valley(VARIABLES.var) ≥ 2
```

Symmetries

- Items of `VARIABLES` can be [reversed](#).
- One and the same constant can be [added](#) to the `var` attribute of all items of `VARIABLES`.

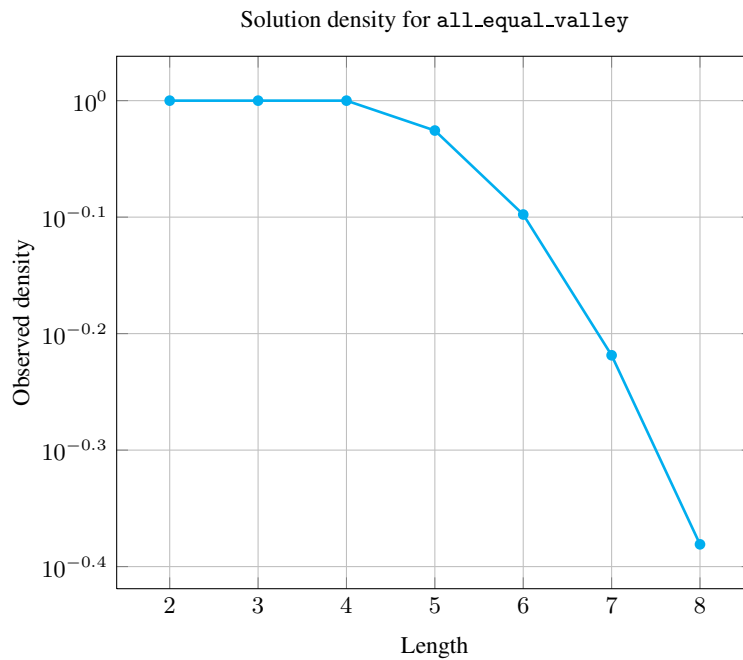
Arg. properties

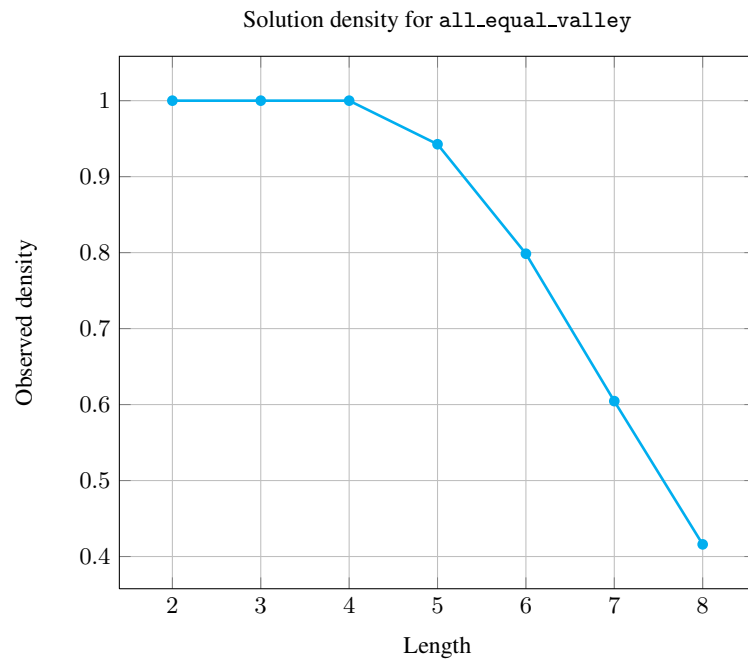
- [Prefix-contractible](#) wrt. `VARIABLES`.
- [Suffix-contractible](#) wrt. `VARIABLES`.

Counting

Length (n)	2	3	4	5	6	7	8
Solutions	9	64	625	7330	93947	1267790	17908059

Number of solutions for `all_equal_valley`: domains $0..n$



**See also**

implied by: `all_equal_valley_min`.

implies: `decreasing_valley`, `increasing_valley`.

related: `all_equal_peak`, `valley`.

Keywords

characteristic of a constraint: `automaton`, `automaton with counters`,
`automaton with same input symbol`.

combinatorial object: `sequence`.

constraint network structure: `sliding cyclic(1) constraint network(2)`.

Cond. implications

- `all_equal_valley(VARIABLES)`
with `valley(VARIABLES.var) > 1`
implies `some_equal(VARIABLES)`.
- `all_equal_valley(VARIABLES)`
with `valley(VARIABLES.var) > 0`
implies `not_all_equal(VARIABLES)`.

Automaton

Figure 5.15 depicts the automaton associated with the `all_equal_valley` constraint. To each pair of consecutive variables (VAR_i, VAR_{i+1}) of the collection `VARIABLES` corresponds a signature variable S_i . The following signature constraint links VAR_i, VAR_{i+1} and S_i : $(VAR_i < VAR_{i+1} \Leftrightarrow S_i = 0) \wedge (VAR_i = VAR_{i+1} \Leftrightarrow S_i = 1) \wedge (VAR_i > VAR_{i+1} \Leftrightarrow S_i = 2)$.

STATES SEMANTICS

s	: initial stationary or increasing mode	$(\{= >\}^*)$
i	: decreasing (before first potential valley) mode	$(< \{< =\}^*)$
j	: increasing (after a valley) mode	$(> \{> =\}^*)$
k	: decreasing (after a valley) mode	$(< \{< =\}^*)$

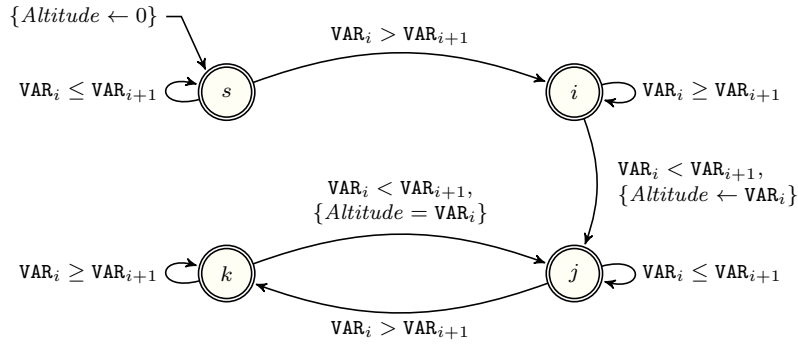


Figure 5.15: Automaton for the `all_equal_valley` constraint (note the conditional transition from state k to state j testing that the counter *Altitude* is equal to VAR_i for enforcing that all valleys are located at the same altitude)

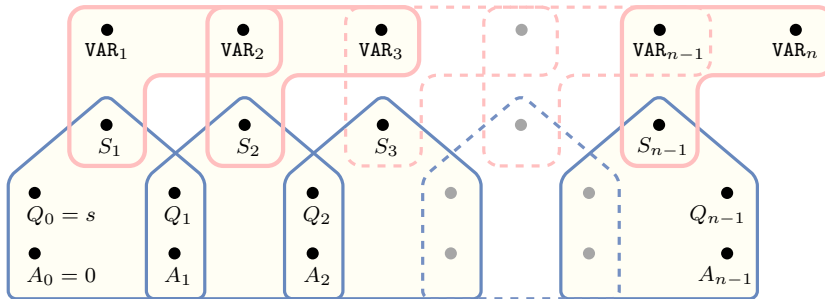


Figure 5.16: Hypergraph of the reformulation corresponding to the automaton of the `all_equal_valley` constraint where A_i stands for the value of the counter *Altitude* (since all states of the automaton are accepting there is no restriction on the last variable Q_{n-1})