

5.36 atleast

	DESCRIPTION	LINKS	GRAPH	AUTOMATON
Origin	CHIP			
Constraint	<code>atleast(N, VARIABLES, VALUE)</code>			
Synonym	<code>count.</code>			
Arguments	N : <code>int</code> VARIABLES : <code>collection(var-dvar)</code> VALUE : <code>int</code>			
Restrictions	$N \geq 0$ $N \leq \text{VARIABLES} $ <code>required(VARIABLES, var)</code>			
Purpose	At least N variables of the VARIABLES collection are assigned value VALUE.			
Example	$(2, \langle 4, 2, 4, 5 \rangle, 4)$ The <code>atleast</code> constraint holds since at least 2 values of the collection $\langle 4, 2, 4, 5 \rangle$ are equal to value 4.			
All solutions	Figure 5.85 gives all solutions to the following non ground instance of the <code>atleast</code> constraint: $V_1 \in [3, 5]$, $V_2 \in [1, 2]$, $V_3 \in [5, 6]$, $V_4 \in [7, 9]$, <code>atleast(2, $\langle V_1, V_2, V_3, V_4 \rangle$, 5)</code> .			
	Figure 5.85: All solutions corresponding to the non ground example of the <code>atleast</code> constraint of the All solutions slot			
Typical	$N > 0$ $N < \text{VARIABLES} $ $ \text{VARIABLES} > 1$			

Symmetries

- Items of VARIABLES are [permutable](#).
- N can be [decreased](#) to any value ≥ 0 .
- An occurrence of a value of VARIABLES.var that is different from VALUE can be [replaced](#) by any other value.

Arg. properties

[Extensible](#) wrt. VARIABLES.

Systems

[occurrenceMin](#) in **Choco**, [count](#) in **Gecode**, [atleast](#) in **Gecode**, [count](#) in **JaCoP**,
[at_least](#) in **MiniZinc**, [count](#) in **SICStus**.

Used in

[alldifferent_except_0](#), [among_diff_0](#), [atmost](#), [global_contiguity](#),
[int_value_precede](#), [ith_pos_different_from_0](#), [minimum_except_0](#),
[nvalues_except_0](#), [period_except_0](#), [sliding_card_skip0](#),
[weighted_partial_alldiff](#).

See also

common keyword: [among](#) (*value constraint*).
comparison swapped: [atmost](#).
implied by: [exactly](#) ($\geq N$ replaced by $= N$).
related: [roots](#).
soft variant: [open_atleast](#) (*open constraint*).

Keywords

characteristic of a constraint: [automaton](#), [automaton with counters](#).
constraint network structure: [alpha-acyclic constraint network\(2\)](#).
constraint type: [value constraint](#).
filtering: [arc-consistency](#).
modelling: [at least](#).

Arc input(s)	VARIABLES
Arc generator	$\text{SELF} \mapsto \text{collection}(\text{variables})$
Arc arity	1
Arc constraint(s)	$\text{variables.var} = \text{VALUE}$
Graph property(ies)	$\text{NARC} \geq N$

Graph model

Since each arc constraint involves only one vertex (**VALUE** is fixed), we employ the *SELF* arc generator in order to produce a graph with a single loop on each vertex.

Parts (A) and (B) of Figure 5.86 respectively show the initial and final graph associated with the **Example** slot. Since we use the **NARC** graph property, the loops of the final graph are stressed in bold.

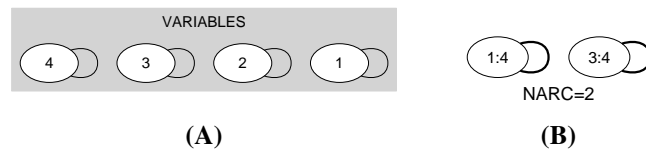


Figure 5.86: Initial and final graph of the at least constraint

Automaton

Figure 5.87 depicts the automaton associated with the `atleast` constraint. To each variable VAR_i of the collection `VARIABLES` corresponds a 0-1 signature variable S_i . The following signature constraint links VAR_i and S_i : $VAR_i = VALUE \Leftrightarrow S_i$. The automaton counts the number of variables of the `VARIABLES` collection that are assigned value `VALUE` and finally checks that this number is greater than or equal to `N`.

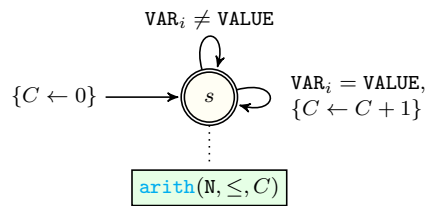


Figure 5.87: Automaton of the `atleast` constraint

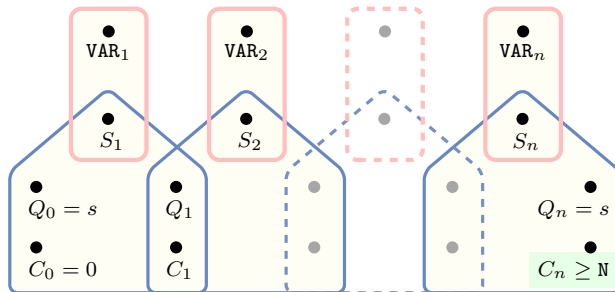


Figure 5.88: Hypergraph of the reformulation corresponding to the automaton (with one counter) of the `atleast` constraint: since all state variables Q_0, Q_1, \dots, Q_n are fixed to the unique state s of the automaton, the transitions constraints share only the counter variable C and the constraint network is Berge-acyclic