AUTOMATON

1400

5.193 inflexion

	DESCRIPTION	LINKS	AUTOMATON
Origin	N. Beldiceanu		
Constraint	<pre>inflexion(N, VARIABLES)</pre>		
Arguments	N : dvar VARIABLES : collection(var-dvar)	
Restrictions	$\begin{split} & \texttt{N} \geq 0 \\ & \texttt{N} \leq \max(0, \texttt{VARIABLES} - 2) \\ & \texttt{required}(\texttt{VARIABLES}, \texttt{var}) \end{split}$		
Purpose	N is equal to the number of times • $X_i \operatorname{CTR} X_{i+1} \wedge X_i \neq X_i$ • $X_{i+1} = X_{i+2} \wedge \cdots \wedge X_j$ • $X_{j-1} \neq X_j \wedge X_{j-1} \neg \operatorname{CT}$	that the following conj +1, $_{j-2} = X_{j-1}$, 'R X_j .	unctions of constraints hold:
	where X_k is the k^{th} item of the V CTR is $<$ or $>$.	VARIABLES collection a	nd $1 \leq i, i+2 \leq j, j \leq n$ and
Example	$\begin{array}{c} (3, \langle 1, 1, 4, 8, 8, 2, 7, 1 \rangle) \\ (0, \langle 1, 1, 4, 4, 6, 6, 7, 9 \rangle) \\ (7, \langle 1, 0, 2, 0, 7, 2, 7, 1, 2 \rangle) \end{array}$		
	The first inflexion constraint three inflexions peaks that respect	holds since the sequentively correspond to value	nce 1 1 4 8 8 2 7 1 contains es 8, 2 and 7.
All solutions	Figure 5.429 gives all solutions to constraint: $\mathbb{N} \in \{0, 2\}, V_1 = $ inflexion($\mathbb{N}, \langle V_1, V_2, V_3, V_4, V_5$	the following non group $V_2 \in [2,3], V_3 \in $	bund instance of the inflexion $[1,2], V_4 \in [1,2], V_5 = 3,$
Typical	$\begin{split} & \texttt{N} > 0 \\ & \texttt{VARIABLES} > 2 \\ & \texttt{range}(\texttt{VARIABLES.var}) > 1 \end{split}$		
Symmetries	 Items of VARIABLES can be One and the same constant VARIABLES. 	be reversed.	e var attribute of all items of
Arg. properties	Functional dependency: N determ	ined by VARIABLES.	
Usage	Useful for constraining the numbe	r of inflexions of a sequ	ence of domain variables.



Figure 5.428: Illustration of the first example of the **Example** slot: a sequence of eight variables V_1 , V_2 , V_3 , V_4 , V_5 , V_6 , V_7 , V_8 respectively fixed to values 1, 1, 4, 8, 8, 2, 7, 1 and its three inflexions in red, two peaks and one valley ($\mathbb{N} = 3$)



Figure 5.429: All solutions corresponding to the non ground example of the inflexion constraint of the **All solutions** slot where each inflexion (i.e. peak or valley) is coloured in orange or cyan

Remark Since the arity of the rently described with

Since the arity of the arc constraint is not fixed, the inflexion constraint cannot be currently described with the graph-based representation. However, this would not hold anymore if we were introducing a slot that specifies how to merge adjacent vertices of the final graph.

Counting

Length (n)	2	3	4	5	6	7	8
Solutions	9	64	625	7776	117649	2097152	43046721
Number of solutions for inflexion: domains $0n$							





Length (n)		2	3	4	5	6	7	8
Total		9	64	625	7776	117649	2097152	43046721
Parameter value	0	9	36	135	498	1841	6856	25731
	1	-	28	320	2588	18494	125284	828120
	2	-	-	170	3348	44058	492320	5069970
	3	-	-	-	1342	40446	778936	12341184
	4	-	-	-	-	12810	549152	14547186
	5	-	-	-	-	-	144604	8354520
	6	-	-	-	-	-	-	1880010

Solution count for inflexion: domains 0..n





constraint arguments: reverse of a constraint, pure functional dependency.

constraint network structure: sliding cyclic(1) constraint network(2).

filtering: glue matrix.

modelling: functional dependency.

- $\label{eq:cond.implications} \begin{array}{ll} \bullet \texttt{inflexion(N,VARIABLES)} \\ & \texttt{with } \texttt{N} > 0 \\ & \texttt{implies atleast_nvalue(NVAL,VARIABLES)} \\ & \texttt{when } \texttt{NVAL} = 2. \end{array}$
 - inflexion(N,VARIABLES) with valley(VARIABLES.var) = 0 implies peak(N,VARIABLES).
 - inflexion(N, VARIABLES) with peak(VARIABLES.var) = 0 implies valley(N, VARIABLES).

See also

Keywords

Automaton

Figure 5.430 depicts the automaton associated with the inflexion constraint. To each pair of consecutive variables (VAR_i, VAR_{i+1}) of the collection VARIABLES corresponds a signature variable S_i . The following signature constraint links VAR_i , VAR_{i+1} and S_i : $(VAR_i < VAR_{i+1} \Leftrightarrow S_i = 0) \land (VAR_i = VAR_{i+1} \Leftrightarrow S_i = 1) \land (VAR_i > VAR_{i+1} \Leftrightarrow S_i = 2)$.



Figure 5.430: Automaton of the inflexion constraint (state *s* means that we are in *stationary* mode, state *i* means that we are in *increasing* mode, state *j* means that we are in *decreasing* mode, a new inflexion is detected each time we switch from increasing to decreasing mode – or conversely from decreasing to increasing mode – and the counter C is incremented accordingly)



Figure 5.431: Hypergraph of the reformulation corresponding to the automaton of the inflexion constraint

	$s (=^*)$	$i (< \{< =\}^*)$	$j (> \{> =\}^*)$
$s\left(=^{*} ight)$	0 —×—	to the	
$i \ (<\{< =\}^*)$		$\overrightarrow{C} + 1 + \overleftarrow{C}$	$\overrightarrow{C} + \overleftarrow{C}$
$j (> \{> =\}^*)$			\overrightarrow{C} + 1 + \overleftarrow{C}

Glue matrix where \overrightarrow{C} and \overleftarrow{C} resp. represent the counter value C at the end of a prefix and at the end of the corresponding reverse suffix that partitions the sequence VARIABLES.

Figure 5.432: Glue matrix associated with the automaton of the inflexion constraint



Figure 5.433: Illustrating the use of the state pair (j, j) of the glue matrix for linking N with the counters variables obtained after reading the prefix 1, 1, 4, 8, 8, 2 and corresponding suffix 2, 7, 1 of the sequence 1, 1, 4, 8, 8, 2, 7, 1; note that the suffix 2, 7, 1 (in pink) is proceed in reverse order; the left (resp. right) table shows the initialisation (for i = 0) and the evolution (for i > 0) of the state of the automaton and its counter C upon reading the prefix 1, 1, 4, 8, 8, 2 (resp. the reverse suffix 1, 7, 2).