

5.224 `lex_chain_greatereq`

	DESCRIPTION	LINKS	GRAPH
Origin	Derived from <code>lex_chain_lesseq</code>		
Constraint	<code>lex_chain_greatereq(VECTORS)</code>		
Usual name	<code>lex_chain</code>		
Type	VECTOR : <code>collection(var-dvar)</code>		
Argument	VECTORS : <code>collection(vec - VECTOR)</code>		
Restrictions	$ \text{VECTOR} \geq 1$ <code>required(VECTOR, var)</code> <code>required(VECTORS, vec)</code> <code>same_size(VECTORS, vec)</code>		
Purpose	<p>For each pair of consecutive vectors VECTOR_i and VECTOR_{i+1} of the <code>VECTORS</code> collection we have that VECTOR_i is lexicographically greater than or equal to VECTOR_{i+1}. Given two vectors, \vec{X} and \vec{Y} of n components, $\langle X_0, \dots, X_{n-1} \rangle$ and $\langle Y_0, \dots, Y_{n-1} \rangle$, \vec{X} is lexicographically greater than or equal to \vec{Y} if and only if $n = 0$ or $X_0 > Y_0$ or $X_0 = Y_0$ and $\langle X_1, \dots, X_{n-1} \rangle$ is lexicographically greater than or equal to $\langle Y_1, \dots, Y_{n-1} \rangle$.</p>		
Example	<code>((vec - <5, 2, 6, 2>, vec - <5, 2, 6, 2>, vec - <5, 2, 3, 9>))</code>		
	<p>The <code>lex_chain_greatereq</code> constraint holds since:</p> <ul style="list-style-type: none"> • The first vector $\langle 5, 2, 6, 2 \rangle$ of the <code>VECTORS</code> collection is lexicographically greater than or equal to the second vector $\langle 5, 2, 6, 2 \rangle$ of the <code>VECTORS</code> collection. • The second vector $\langle 5, 2, 6, 2 \rangle$ of the <code>VECTORS</code> collection is lexicographically greater than or equal to the third vector $\langle 5, 2, 3, 9 \rangle$ of the <code>VECTORS</code> collection. 		
Typical	$ \text{VECTOR} > 1$ $ \text{VECTORS} > 1$		
Arg. properties	<ul style="list-style-type: none"> • Contractible wrt. <code>VECTORS</code>. • Suffix-contractible wrt. <code>VECTORS.vec</code> (<i>remove items from same position</i>). 		
Usage	<p>This constraint was motivated for breaking symmetry: more precisely when one wants to lexicographically order the consecutive columns of a matrix of decision variables. A further motivation is that using a set of lexicographic ordering constraints between two vectors does usually not allow to come up with a complete pruning.</p>		

- Algorithm** A filtering algorithm achieving [arc-consistency](#) for a chain of lexicographical ordering constraints is presented in [95].
- Six different ways of integrating a chain of lexicographical ordering constraints within non-overlapping constraints like [diffn](#) or [geost](#) and within their corresponding necessary condition like the [cumulative](#) constraint are shown in [3].
- See also** **common keyword:** [lex_between](#), [lex_greater](#), [lex_less](#), [lex_lesseq](#) (*lexicographic order*).
- implied by:** [lex_chain_greater](#) (*non-strict order implied by strict order*).
- part of system of constraints:** [lex_greatereq](#).
- used in graph description:** [lex_greatereq](#).
- Keywords** **characteristic of a constraint:** [vector](#).
- constraint type:** [system of constraints](#), [decomposition](#), [order constraint](#).
- filtering:** [arc-consistency](#).
- heuristics:** [heuristics and lexicographical ordering](#).
- symmetry:** [symmetry](#), [matrix symmetry](#), [lexicographic order](#).

Arc input(s)	VECTORS
Arc generator	$\text{PATH} \mapsto \text{collection}(\text{vectors1}, \text{vectors2})$
Arc arity	2
Arc constraint(s)	$\text{lex_lesseq}(\text{vectors1.vec}, \text{vectors2.vec})$
Graph property(ies)	$\text{NARC} = \text{VECTORS} - 1$

Graph model

Parts (A) and (B) of Figure 5.479 respectively show the initial and final graph associated with the **Example** slot. Since we use the $\overline{\text{NARC}}$ graph property, the arcs of the final graph are stressed in bold. The $\text{lex_chain_greatereq}$ constraint holds since all the arc constraints of the initial graph are satisfied.

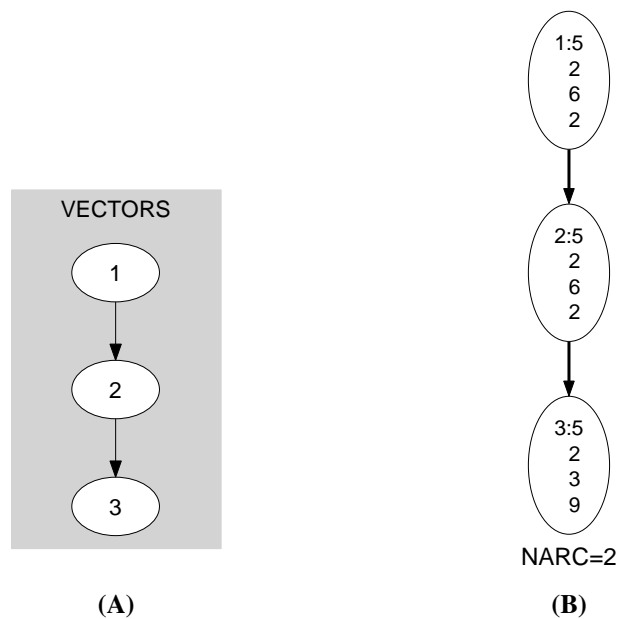


Figure 5.479: Initial and final graph of the $\text{lex_chain_greatereq}$ constraint

Signature

Since we use the PATH arc generator on the VECTORS collection the number of arcs of the initial graph is equal to $|\text{VECTORS}| - 1$. For this reason we can rewrite $\text{NARC} = |\text{VECTORS}| - 1$ to $\text{NARC} \geq |\text{VECTORS}| - 1$ and simplify $\overline{\text{NARC}}$ to $\overline{\text{NARC}}$.

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