5.236 longest_decreasing_sequence

	DESCRIPTION	LINKS	AUTOMATON				
Origin	constraint on sequences						
Constraint	$\verb+longest_decreasing_sequence(L, VARIABLES)$						
Synonym	<pre>size_longest_decreasing_sequence.</pre>						
Arguments	L : dvar VARIABLES : collection(var-dvar)						
Restrictions	$\label{eq:lass} \begin{array}{l} \mathtt{L} \geq 0 \\ \mathtt{L} < \mathtt{range}(\mathtt{VARIABLES}.\mathtt{var}) \\ \mathtt{required}(\mathtt{VARIABLES},\mathtt{var}) \end{array}$						
Purpose	L is the largest difference betwee sequences of the collection VAI A sequence of consecutive van of the collection of variables V following conditions simultance • $X_i \ge X_{i+1} \ge \cdots \ge X$ • $i = 1$ or $X_{i-1} < X_i$, • $i = VARIABLES $ or X_j	the first and the last of RIABLES. Triables $X_i, X_{i+1}, \ldots, X_i$ VARIABLES is a maximum cously apply: T_j , $< X_{j+1}$.	value of the maximum decreasing C_j $(1 \le i \le j \le VARIABLES)$ <i>im decreasing sequence</i> if all the				
Example	$(0, \langle 0, 1, 2, 5 \rangle)$ $(0, \langle 8, 8 \rangle)$ $(6, \langle 10, 8, 8, 6, 4, 9, 10, 8 \rangle)$ Figure 5.504 gives a graphical slot with its two maximum dec corresponding longest_decrease L is fixed to the maximum size of the statement of the maximum size of the statement of the st	l representation of the reasing sequences in re asing_sequence constr 3.	third example of the Example d of respective size 6 and 2. The raint holds since its first argument				
Typical	$\begin{array}{l} {\tt L} > 0 \\ {\tt VARIABLES} > 1 \\ {\tt nval}({\tt VARIABLES.var}) > 2 \end{array}$						
Symmetry	One and the same constant can	be added to the var attr	ibute of all items of VARIABLES.				
Arg. properties	Functional dependency: L determined by VARIABLES.						
Counting							

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Figure 5.504: Illustration of the third example of the **Example** slot: a sequence of eight variables V_1 , V_2 , V_3 , V_4 , V_5 , V_6 , V_7 , V_8 respectively fixed to values 10, 8, 8, 6, 4, 9, 10, 8 and its two maximum decreasing sequences in red of respective size 10 - 4 = 6 and 10 - 8 = 2.

Length (n)	2	3	4	5	6	7	8
Solutions	9	64	625	7776	117649	2097152	43046721
Number of solutions for longest_decreasing_sequence: domains $0n$							



Solution density for longest_decreasing_sequence



Length (n)		2	3	4	5	6	7	8
Total		9	64	625	7776	117649	2097152	43046721
Parameter 4	0	6	20	70	252	924	3432	12870
	1	2	18	122	750	4412	25382	144314
	2	1	16	161	1398	11361	89132	685090
	3	-	10	162	1942	20816	211106	2074365
	4	-	-	110	2024	28930	375084	4603682
value	5	-	-	-	1410	30134	506766	7792840
6 7 8	6	-	-	-	-	21072	522648	10197174
	7	-	-	-	-	-	363602	10379696
	8	-	-	-	-	-	-	7156690

Solution count for longest_decreasing_sequence: domains 0..n



 $Solution \ {\tt density} \ for \ {\tt longest_decreasing_sequence}$







longest_increasing_sequence,

AUTOMATON

 Keywords
 characteristic of a constraint: automaton, automaton with counters, automaton with same input symbol.
 automaton with counters, automaton, automaton with counters, automaton with same input symbol.

 combinatorial object: sequence.
 constraint arguments: reverse of a constraint, pure functional dependency.

 filtering: glue matrix.
 modelling: functional dependency.

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Automaton

Figure 5.505 depicts the automaton associated with the <code>longest_decreasing_sequence</code> constraint.

STATES SEMANTICS

$$s :: increasing mode (\{<|=\}^*) t :: decreasing mode (\{<|=\}^*) VAR_i > VAR_i - VAR_{i+1}, C \leftarrow VAR_i - VAR_{i+1}, C \leftarrow VAR_i - VAR_{i+1}, C \leftarrow VAR_i - VAR_{i+1}, VAR_i \leq VAR_{i+1}, VAR_i \leq VAR_{i+1}, VAR_i < VAR_{i+1}, C \leftarrow C + VAR_i - VAR_{i+1}, C \leftarrow C + VAR_i - VAR_{i+1}, C \leftarrow C + VAR_i - VAR_{i+1}, S (\{<|=\}^*) t (<\{<|=\}^*) max(\vec{M}, \vec{M}) max(\vec{M}, \vec{M})$$

end of a prefix and at the end of the corresponding reverse suffix that partitions the sequence VARIABLES. $t (> \{> | =$

	$s(\{ \ge - \})$	
$ = \}^*)$	$\max(\overrightarrow{M},\overleftarrow{M})$	$\max(\overrightarrow{M}, \overleftarrow{M})$
=}*)	$\max(\overrightarrow{M},\overleftarrow{M})$	$\max(\overrightarrow{M},\overrightarrow{C}+\overleftarrow{C},\overleftarrow{M})$

Figure 5.505: Automaton of the longest_decreasing_sequence constraint and its glue matrix (note that the reverse of the longest_decreasing_sequence constraint is the longest_increasing_sequence constraint)



Figure 5.506: Hypergraph of the reformulation corresponding to the automaton of the longest_decreasing_sequence constraint