

5.262 minimum

	DESCRIPTION	LINKS	GRAPH	AUTOMATON
Origin	CHIP			
Constraint	<code>minimum(MIN, VARIABLES)</code>			
Synonym	<code>min.</code>			
Arguments	<pre> MIN : <code>dvar</code> VARIABLES : <code>collection(var-dvar)</code> </pre>			
Restrictions	<pre> VARIABLES > 0 <code>required(VARIABLES, var)</code> </pre>			
Purpose	MIN is the minimum value of the collection of domain variables VARIABLES.			
Example	<pre> (2, <3, 2, 7, 2, 6>) (7, <8, 8, 7, 8, 7>) </pre> <p>The first minimum constraint holds since its first argument $MIN = 2$ is set to the minimum value of the collection $\langle 3, 2, 7, 2, 6 \rangle$.</p>			
Typical	<pre> VARIABLES > 1 <code>range(VARIABLES.var) > 1</code> </pre>			
Symmetries	<ul style="list-style-type: none"> Items of VARIABLES are <i>permutable</i>. All occurrences of two distinct values of VARIABLES.var can be <i>swapped</i>. One and the same constant can be <i>added</i> to MIN as well as to the var attribute of all items of VARIABLES. 			
Arg. properties	<ul style="list-style-type: none"> Functional dependency: MIN determined by VARIABLES. Aggregate: <code>MIN(min), VARIABLES(union)</code>. 			
Usage	In some project scheduling problems one has to introduce dummy activities that correspond for instance to the starting time of a given set of activities. In this context one can use the minimum constraint to get the minimum starting time of a set of tasks.			
Remark	<p>Note that <code>minimum</code> is a constraint and not just a function that computes the minimum value of a collection of variables: potential values of MIN influence the variables of VARIABLES, and reciprocally potential values that can be assigned to variables of VARIABLES influence MIN.</p> <p>The minimum constraint is called <code>min</code> in JaCoP (http://www.jacop.eu/).</p>			

Algorithm

A filtering algorithm for the minimum constraint is described in [27].

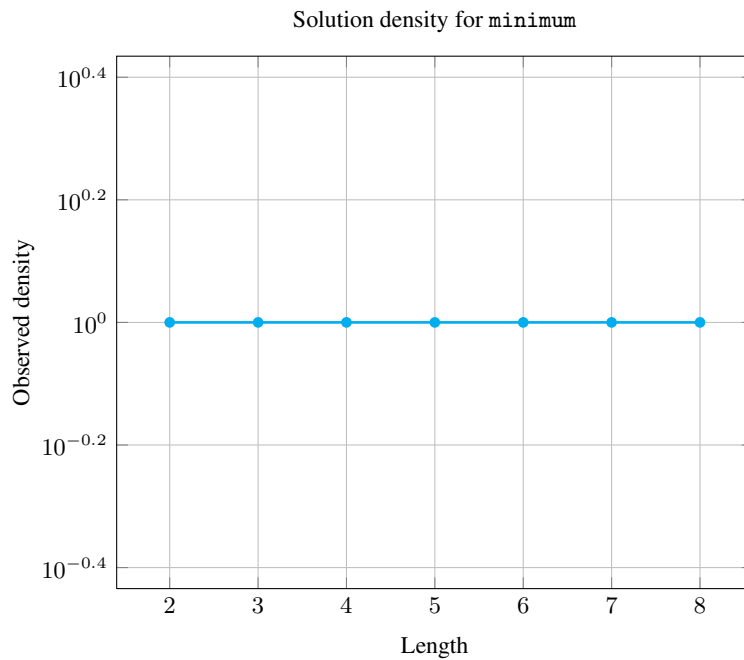
The minimum constraint is **entailed** if all the following conditions hold:

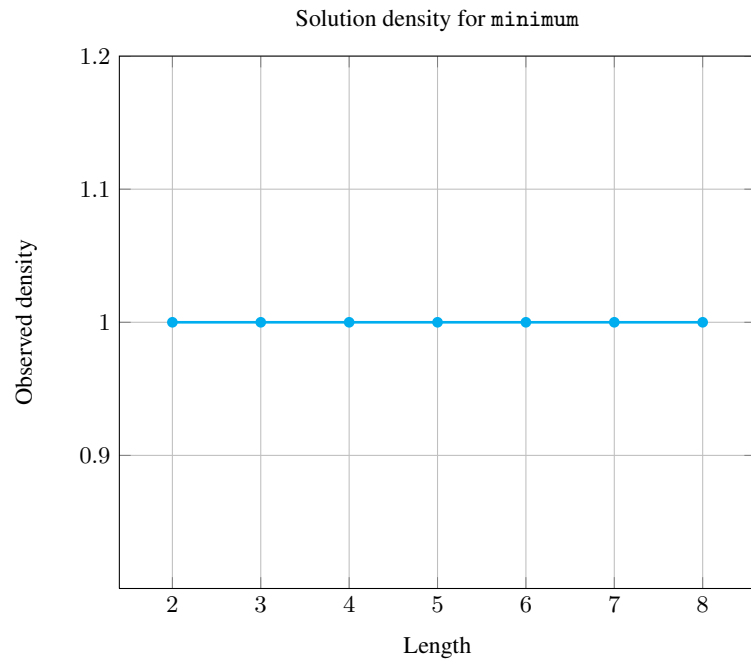
1. MIN is fixed.
2. At least one variable of VARIABLES is assigned value MIN.
3. All variables of VARIABLES have their minimum value greater than or equal to value MIN.

Counting

Length (n)	2	3	4	5	6	7	8
Solutions	9	64	625	7776	117649	2097152	43046721

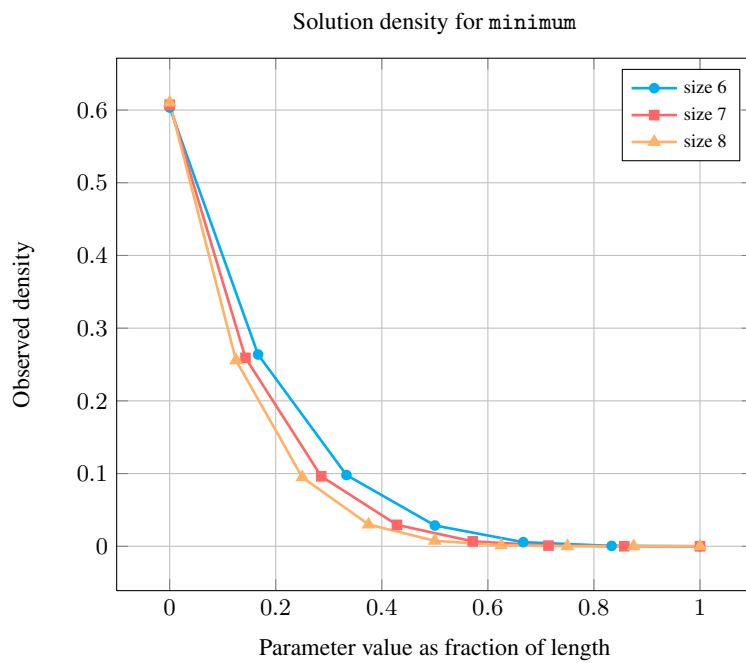
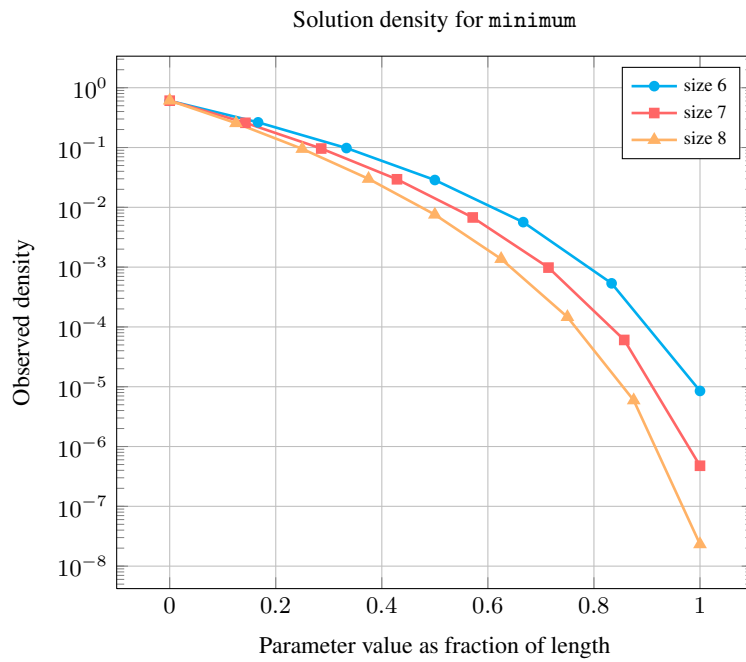
Number of solutions for minimum: domains $0..n$





Length (n)		2	3	4	5	6	7	8
Total		9	64	625	7776	117649	2097152	43046721
Parameter value	0	5	37	369	4651	70993	1273609	26269505
	1	3	19	175	2101	31031	543607	11012415
	2	1	7	65	781	11529	201811	4085185
	3	-	1	15	211	3367	61741	1288991
	4	-	-	1	31	665	14197	325089
	5	-	-	-	1	63	2059	58975
	6	-	-	-	-	1	127	6305
	7	-	-	-	-	-	1	255
	8	-	-	-	-	-	-	1

Solution count for minimum: domains $0..n$



Systems

[min](#) in [Choco](#), [min](#) in [Gecode](#), [min](#) in [JaCoP](#), [minimum](#) in [MiniZinc](#), [minimum](#) in [SICStus](#).

Used in	<code>minimum_greater_than</code> , <code>next_element</code> , <code>next_greater_element</code> .
See also	<p>common keyword: <code>maximum</code> (<i>order constraint</i>).</p> <p>comparison swapped: <code>maximum</code>.</p> <p>generalisation: <code>minimum_modulo</code> (<i>variable replaced by variable mod constant</i>).</p> <p>implied by: <code>and</code>.</p> <p>implies: <code>between_min_max</code>, <code>in</code>.</p> <p>soft variant: <code>minimum_except_0</code> (<i>value 0 is ignored</i>), <code>open_minimum</code> (<i>open constraint</i>).</p> <p>specialisation: <code>min_n</code> (<i>minimum or order n replaced by absolute minimum</i>).</p> <p>uses in its reformulation: <code>cycle</code>.</p>
Keywords	<p>characteristic of a constraint: <code>minimum</code>, <code>maxint</code>, <code>automaton</code>, <code>automaton without counters</code>, <code>reified automaton constraint</code>.</p> <p>constraint arguments: reverse of a constraint, pure functional dependency.</p> <p>constraint network structure: centered cyclic(1) constraint network(1).</p> <p>constraint type: order constraint.</p> <p>filtering: glue matrix, arc-consistency, entailment.</p> <p>modelling: functional dependency.</p>
Cond. implications	<pre> minimum(MIN, VARIABLES) with first(VARIABLES.var) > MIN and last(VARIABLES.var) > MIN implies deepest_valley(DEPTH, VARIABLES). </pre>

Arc input(s)	VARIABLES
Arc generator	<code>CLIQUE</code> \mapsto <code>collection</code> (variables1, variables2)
Arc arity	2
Arc constraint(s)	$\bigvee \left(\begin{array}{l} \text{variables1.key} = \text{variables2.key}, \\ \text{variables1.var} < \text{variables2.var} \end{array} \right)$
Graph property(ies)	<code>ORDER</code> (0, MAXINT, var) = MIN

Graph model

The condition `variables1.key = variables2.key` holds if and only if `variables1` and `variables2` corresponds to the same vertex. It is used in order to enforce to keep all the vertices of the initial graph. `ORDER`(0, MAXINT, var) refers to the source vertices of the graph, i.e., those vertices that do not have any predecessor.

Parts (A) and (B) of Figure 5.555 respectively show the initial and final graph associated with the first example of the **Example** slot. Since we use the `ORDER` graph property, the vertices of rank 0 (without considering the loops) of the final graph are outlined with a thick circle.

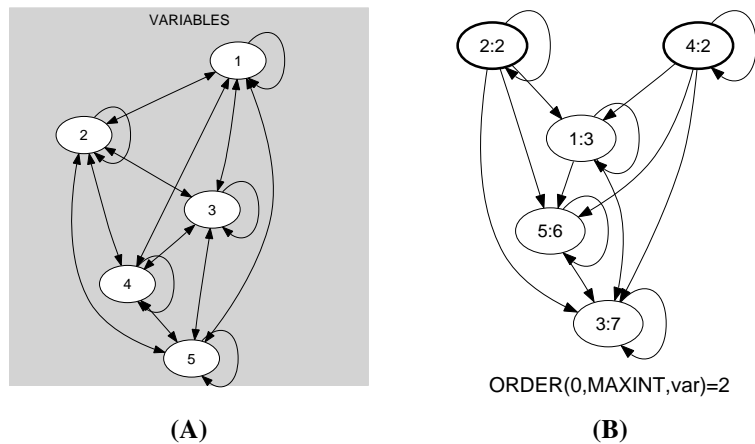


Figure 5.555: Initial and final graph of the minimum constraint

Automaton

Figure 5.556 depicts a first counter free deterministic automaton associated with the minimum constraint. Let VAR_i be the i^{th} variable of the VARIABLES collection. To each pair (MIN, VAR_i) corresponds a signature variable S_i as well as the following signature constraint: $(MIN < VAR_i \Leftrightarrow S_i = 0) \wedge (MIN = VAR_i \Leftrightarrow S_i = 1) \wedge (MIN > VAR_i \Leftrightarrow S_i = 2)$.

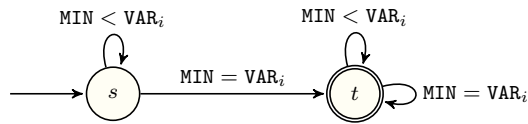


Figure 5.556: Counter free automaton of the minimum constraint

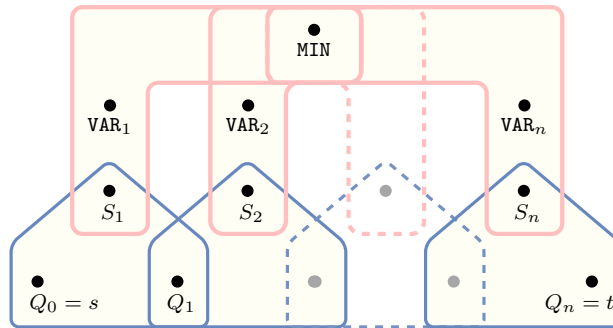


Figure 5.557: Hypergraph of the reformulation corresponding to the automaton of the minimum constraint

Figure 5.557 depicts a second counter free non deterministic automaton associated with the minimum constraint, where the argument MIN is also part of the sequence passed to the automaton.

Figure 5.560 depicts a third deterministic automaton with one counter associated with the minimum constraint, where the argument MIN is unified to the final value of the counter.

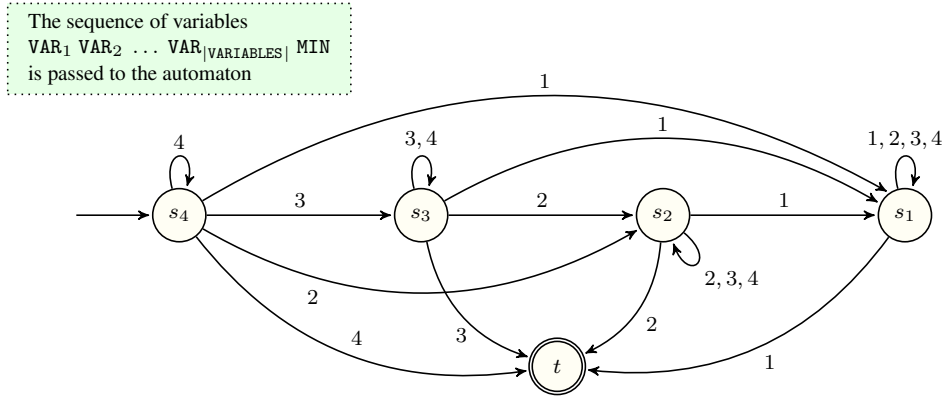


Figure 5.558: Counter free non deterministic automaton of the minimum(MIN, VARIABLES) constraint assuming that the union of the domain of the variables is the set {1, 2, 3, 4} and that the elements of VARIABLES are first passed to the automaton followed by MIN (state s_i means that no value strictly less than value i was found and that value i was already encountered at least once)

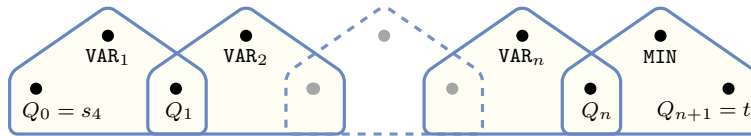


Figure 5.559: Hypergraph of the reformulation corresponding to the counter free non deterministic automaton of the minimum constraint

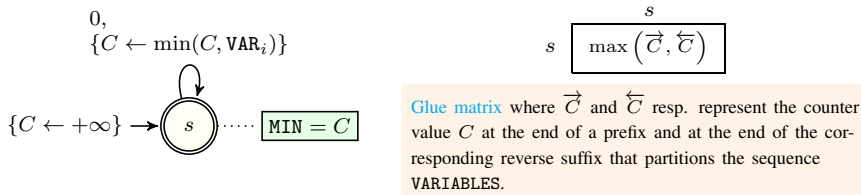


Figure 5.560: Automaton (with one counter) of the minimum constraint and its glue matrix

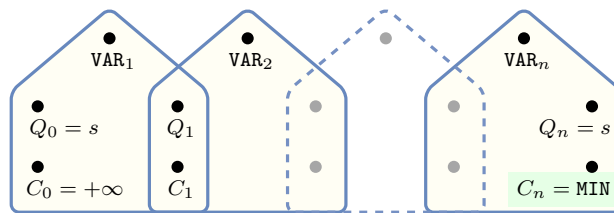


Figure 5.561: Hypergraph of the reformulation corresponding to the automaton (with one counter) of the minimum constraint: since all states variables Q_0, Q_1, \dots, Q_n are fixed to the unique state s of the automaton, the transitions constraints share only the counter variable C and the constraint network is Berge-acyclic

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