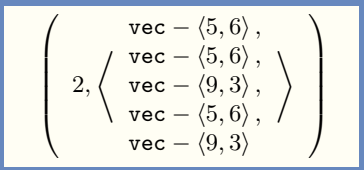


5.290 nvector

	DESCRIPTION	LINKS	GRAPH
Origin	Introduced by G. Chabert as a generalisation of <code>nvalue</code>		
Constraint	<code>nvector(NVEC, VECTORS)</code>		
Synonyms	<code>nvectors</code> , <code>npoint</code> , <code>npoints</code> .		
Type	VECTOR : <code>collection(var-dvar)</code>		
Arguments	NVEC : <code>dvar</code> VECTORS : <code>collection(vec - VECTOR)</code>		
Restrictions	$ \text{VECTOR} \geq 1$ $\text{NVEC} \geq \min(1, \text{VECTORS})$ $\text{NVEC} \leq \text{VECTORS} $ <code>required(VECTORS, vec)</code> <code>same_size(VECTORS, vec)</code>		
Purpose	<p>NVEC is the number of distinct tuples of values taken by the vectors of the collection VECTORS. Two tuples of values $\langle A_1, A_2, \dots, A_m \rangle$ and $\langle B_1, B_2, \dots, B_m \rangle$ are <i>distinct</i> if and only if there exist an integer $i \in [1, m]$ such that $A_i \neq B_i$.</p>		
Example			
	<p>The <code>nvector</code> constraint holds since its first argument $\text{NVEC} = 2$ is set to the number of distinct tuples of values (i.e., tuples $\langle 5, 6 \rangle$ and $\langle 9, 3 \rangle$) occurring within the collection VECTORS. Figure 5.604 depicts with a thick rectangle a possible initial domain for each of the five vectors and with a grey circle each tuple of values of the corresponding solution.</p>		
Typical	$ \text{VECTOR} > 1$ $\text{NVEC} > 1$ $\text{NVEC} < \text{VECTORS} $ $ \text{VECTORS} > 1$		
Symmetries	<ul style="list-style-type: none"> • Items of VECTORS are <code>permutable</code>. • Items of VECTORS.vec are <code>permutable</code> (<i>same permutation used</i>). • All occurrences of two distinct tuples of values of VECTORS.vec can be <code>swapped</code>; all occurrences of a tuple of values of VECTORS.vec can be <code>renamed</code> to any unused tuple of values. 		

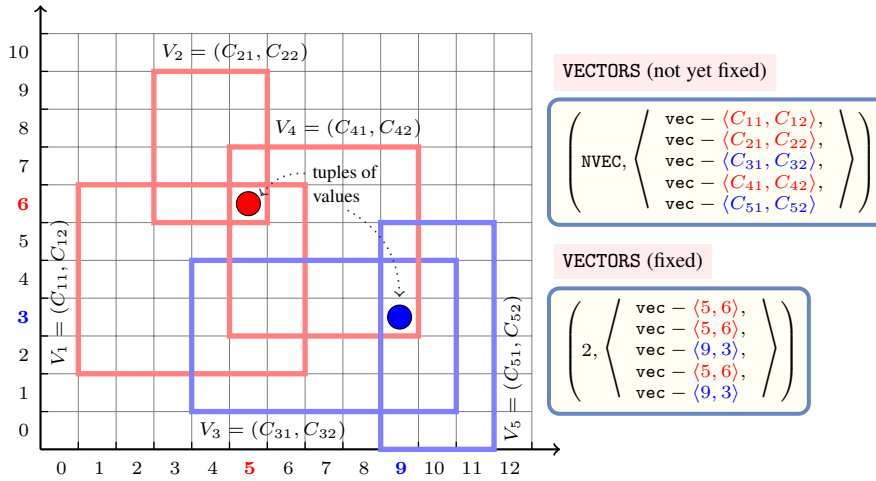


Figure 5.604: Possible initial domains ($C_{11} \in [1, 6], C_{12} \in [2, 6], C_{21} \in [3, 5], C_{22} \in [6, 9], C_{31} \in [4, 10], C_{32} \in [1, 4], C_{41} \in [5, 9], C_{42} \in [3, 7], C_{51} \in [9, 11], C_{52} \in [0, 5]$) and solution corresponding to the **Example** slot: we have two distinct vectors (NVEC = 2)

Arg. properties

- **Functional dependency**: NVEC determined by VECTORS.
- **Contractible** wrt. VECTORS when NVEC = 1 and |VECTORS| > 0.
- **Contractible** wrt. VECTORS when NVEC = |VECTORS|.

Remark

It was shown in [109, 108] that, finding out whether a nvector constraint has a solution or not is NP-hard (i.e., the restriction to the rectangle case and to the atmost side of the nvector were considered for this purpose). This was achieved by reduction from the [rectangle clique partition](#) problem.

Reformulation

Assume the collection VECTORS is not empty (otherwise NVEC = 0). In this context, let n and m respectively denote the number of vectors of the collection VECTORS and the number of components of each vector. Furthermore, let $\alpha_i = \min(C_{1i}, C_{2i}, \dots, C_{ni}), \beta_i = \max(\overline{C_{1i}}, \overline{C_{2i}}, \dots, \overline{C_{ni}}), \gamma_i = \beta_i - \alpha_i + 1, (i \in [1, m])$. By associating to each vector

$$\langle C_{k1}, C_{k2}, \dots, C_{km} \rangle, (k \in [1, n])$$

a variable

$$D_k = \sum_{1 \leq i \leq m} \left(\left(\prod_{i < j \leq m} \gamma_j \right) \cdot (C_{ki} - \alpha_i) \right),$$

the constraint

$$\text{nvector}(\text{NVEC}, \langle \text{vec} - \langle C_{11}, C_{12}, \dots, C_{1m} \rangle, \text{vec} - \langle C_{21}, C_{22}, \dots, C_{2m} \rangle, \dots, \text{vec} - \langle C_{n1}, C_{n2}, \dots, C_{nm} \rangle \rangle)$$

can be expressed in term of the constraint
`nvalue(NVEC, (D1, D2, ..., Dn))`.

Note that the previous reformulation does not work anymore if the variables have a continuous domain, or if an overflow occurs while propagating the equality constraint $D_k = \sum_{1 \leq i \leq m} \left(\left(\prod_{i < j \leq m} \gamma_j \right) \cdot (C_{ki} - \alpha_i) \right)$ (i.e., the number of components m is too big).

When using this reformulation with respect to the **Example** slot we first introduce $D_1 = 1 \cdot 6 - 3 + (4 \cdot 5 - 20) = 3$, $D_2 = 1 \cdot 6 - 3 + (4 \cdot 5 - 20) = 3$, $D_3 = 1 \cdot 3 - 3 + (4 \cdot 9 - 20) = 16$, $D_4 = 1 \cdot 6 - 3 + (4 \cdot 5 - 20) = 3$, $D_5 = 1 \cdot 3 - 3 + (4 \cdot 9 - 20) = 16$ and then get the constraint `nvalue(2, (3, 3, 16, 3, 16))`.

See also

common keyword: `lex_equal`, `ordered_atleast_nvector`, `ordered_atmost_nvector` (`vector`).

generalisation: `nvectors` (replace an equality with the number of distinct vectors by a comparison with the number of distinct `nvectors`).

implied by: `ordered_nvector`.

implies: `atleast_nvector` (= NVEC replaced by \geq NVEC), `atmost_nvector` (= NVEC replaced by \leq NVEC).

specialisation: `nvalue` (`vector` replaced by `variable`).

Keywords

application area: SLAM problem.

characteristic of a constraint: `vector`.

complexity: rectangle clique partition.

constraint arguments: pure functional dependency.

constraint type: counting constraint, value partitioning constraint.

final graph structure: strongly connected component, equivalence.

modelling: number of distinct equivalence classes, functional dependency.

problems: domination.

Arc input(s)	VECTORS
Arc generator	<code>CLIQUE</code> \mapsto <code>collection</code> (vectors1, vectors2)
Arc arity	2
Arc constraint(s)	<code>lex_equal</code> (vectors1.vec, vectors2.vec)
Graph property(ies)	<code>NSCC</code> = NVEC
Graph class	<code>EQUIVALENCE</code>

Graph model

Parts (A) and (B) of Figure 5.605 respectively show the initial and final graph associated with the **Example** slot. Since we use the `NSCC` graph property we show the different strongly connected components of the final graph. Each strongly connected component corresponds to a tuple of values that is assigned to some vectors of the `VECTORS` collection. The 2 following tuple of values $\langle 5, 6 \rangle$ and $\langle 9, 3 \rangle$ are used by the vectors of the `VECTORS` collection.

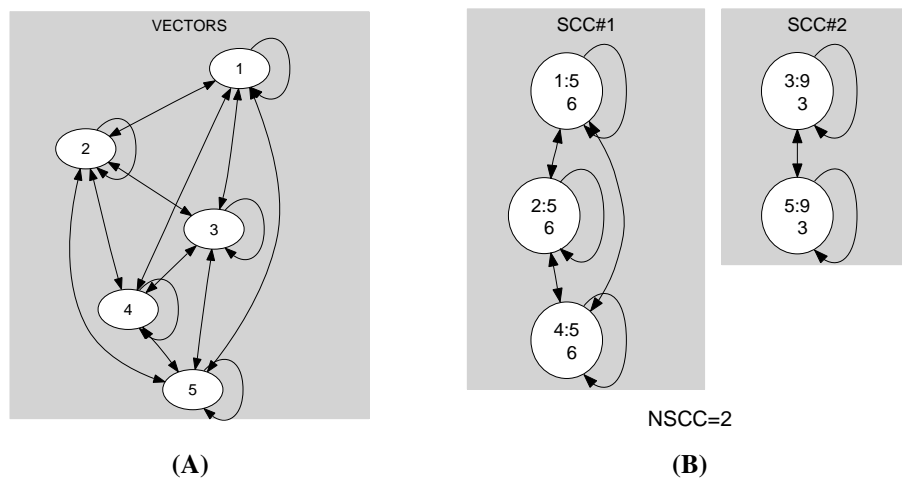


Figure 5.605: Initial and final graph of the nvector constraint