

5.398 symmetric_cardinality

	DESCRIPTION	LINKS	GRAPH
Origin	Derived from <code>global_cardinality</code> by W. Kocjan.		
Constraint	<code>symmetric_cardinality(VARS, VALS)</code>		
Arguments	VARS : <code>collection(idvar-int, var-svar, l-int, u-int)</code> VALS : <code>collection(idval-int, val-svar, l-int, u-int)</code>		
Restrictions	<pre> required(VARS, [idvar, var, l, u]) VARS ≥ 1 VARS.idvar ≥ 1 VARS.idvar ≤ VARS distinct(VARS, idvar) VARS.l ≥ 0 VARS.l ≤ VARS.u VARS.u ≤ VALS required(VALS, [idval, val, l, u]) VALS ≥ 1 VALS.idval ≥ 1 VALS.idval ≤ VALS distinct(VALS, idval) VALS.l ≥ 0 VALS.l ≤ VALS.u VALS.u ≤ VARS </pre>		

Purpose

Put in relation two sets: for each element of one set gives the corresponding elements of the other set to which it is associated. In addition, it constraints the number of elements associated with each element to be in a given interval.

Example

$$\left(\begin{array}{l} \left\langle \begin{array}{llll} \text{idvar} - 1 & \text{var} - \{3\} & 1 - 0 & \text{u} - 1, \\ \text{idvar} - 2 & \text{var} - \{1\} & 1 - 1 & \text{u} - 2, \\ \text{idvar} - 3 & \text{var} - \{1, 2\} & 1 - 1 & \text{u} - 2, \\ \text{idvar} - 4 & \text{var} - \{1, 3\} & 1 - 2 & \text{u} - 3 \end{array} \right\rangle, \\ \left\langle \begin{array}{llll} \text{idval} - 1 & \text{val} - \{2, 3, 4\} & 1 - 3 & \text{u} - 4, \\ \text{idval} - 2 & \text{val} - \{3\} & 1 - 1 & \text{u} - 1, \\ \text{idval} - 3 & \text{val} - \{1, 4\} & 1 - 1 & \text{u} - 2, \\ \text{idval} - 4 & \text{val} - \emptyset & 1 - 0 & \text{u} - 1 \end{array} \right\rangle \end{array} \right)$$

The `symmetric_cardinality` constraint holds since:

- $3 \in \text{VARS}[1].\text{var} \Leftrightarrow 1 \in \text{VALS}[3].\text{val}$,
- $1 \in \text{VARS}[2].\text{var} \Leftrightarrow 2 \in \text{VALS}[1].\text{val}$,
- $1 \in \text{VARS}[3].\text{var} \Leftrightarrow 3 \in \text{VALS}[1].\text{val}$,
- $2 \in \text{VARS}[3].\text{var} \Leftrightarrow 3 \in \text{VALS}[2].\text{val}$,

- $1 \in \text{VARS}[4].\text{var} \Leftrightarrow 4 \in \text{VALS}[1].\text{val}$,
- $3 \in \text{VARS}[4].\text{var} \Leftrightarrow 4 \in \text{VALS}[3].\text{val}$,
- The number of elements of $\text{VARS}[1].\text{var} = \{3\}$ belongs to interval $[0, 1]$,
- The number of elements of $\text{VARS}[2].\text{var} = \{1\}$ belongs to interval $[1, 2]$,
- The number of elements of $\text{VARS}[3].\text{var} = \{1, 2\}$ belongs to interval $[1, 2]$,
- The number of elements of $\text{VARS}[4].\text{var} = \{1, 3\}$ belongs to interval $[2, 3]$,
- The number of elements of $\text{VALS}[1].\text{val} = \{2, 3, 4\}$ belongs to interval $[3, 4]$,
- The number of elements of $\text{VALS}[2].\text{val} = \{3\}$ belongs to interval $[1, 1]$,
- The number of elements of $\text{VALS}[3].\text{val} = \{1, 4\}$ belongs to interval $[1, 2]$,
- The number of elements of $\text{VALS}[4].\text{val} = \emptyset$ belongs to interval $[0, 1]$.

Typical

$|\text{VARS}| > 1$
 $|\text{VALS}| > 1$

Symmetries

- Items of VARS are [permutable](#).
- Items of VALS are [permutable](#).

Usage

The most simple example of applying `symmetric_gcc` is a variant of personnel [assignment](#) problem, where one person can be assigned to perform between n and m ($n \leq m$) jobs, and every job requires between p and q ($p \leq q$) persons. In addition every job requires different kind of skills. The previous problem can be modelled as follows:

- For each person we create an item of the VARS collection,
- For each job we create an item of the VALS collection,
- There is an arc between a person and the particular job if this person is qualified to perform it.

Remark

The `symmetric_gcc` constraint generalises the [global_cardinality](#) constraint by allowing a variable to take more than one value.

Algorithm

A first [flow](#)-based [arc-consistency](#) algorithm for the `symmetric_cardinality` constraint is described in [241]. A second [arc-consistency](#) filtering algorithm exploiting matching theory [148] is described in [129, 130].

See also

common keyword: [link_set_to_booleans](#) (*constraint involving set variables*).

generalisation: [symmetric_gcc](#) (*fixed interval replaced by variable*).

root concept: [global_cardinality](#).

used in graph description: [in_set](#).

Keywords

application area: [assignment](#).

combinatorial object: [relation](#).

constraint arguments: [constraint involving set variables](#).

constraint type: [decomposition](#), [timetabling constraint](#).

filtering: [flow](#), [bipartite matching](#).

Arc input(s)	VARS VALS
Arc generator	<i>PRODUCT</i> \mapsto <code>collection(vars, vals)</code>
Arc arity	2
Arc constraint(s)	<ul style="list-style-type: none"> • <code>in_set(vars.idvar, vals.val) \Leftrightarrow in_set(vals.idval, vars.var)</code> • <code>vars.l \leq card_set(vars.var)</code> • <code>vars.u \geq card_set(vars.var)</code> • <code>vals.l \leq card_set(vals.val)</code> • <code>vals.u \geq card_set(vals.val)</code>
Graph property(ies)	<u>NARC</u> = $ VARS * VALS $

Graph model

The graph model used for the `symmetric_cardinality` is similar to the one used in the `domain_constraint` or in the `link_set_to_booleans` constraints: we use an equivalence in the arc constraint and ask all arc constraints to hold.

Parts (A) and (B) of Figure 5.756 respectively show the initial and final graph associated with the **Example** slot. Since we use the NARC graph property, all the arcs of the final graph are stressed in bold.

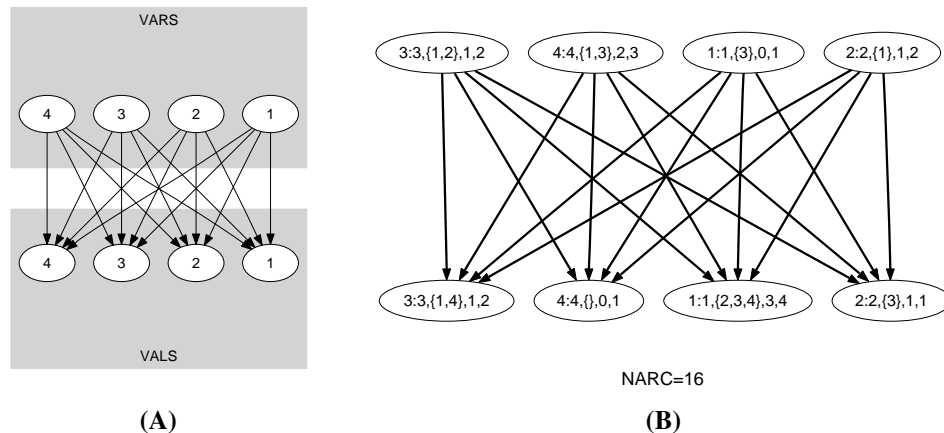


Figure 5.756: Initial and final graph of the `symmetric_cardinality` constraint

Signature

Since we use the *PRODUCT* arc generator on the collections VARS and VALS, the number of arcs of the initial graph is equal to $|VARS| \cdot |VALS|$. Therefore the maximum number of arcs of the final graph is also equal to $|VARS| \cdot |VALS|$ and we can rewrite $NARC = |VARS| \cdot |VALS|$ to $NARC \geq |VARS| \cdot |VALS|$. So we can simplify NARC to \overline{NARC} .

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