

## 5.413 used\_by\_interval

	DESCRIPTION	LINKS	GRAPH
<b>Origin</b>	Derived from <code>used_by</code> .		
<b>Constraint</b>	<code>used_by_interval(VARIABLES1, VARIABLES2, SIZE_INTERVAL)</code>		
<b>Arguments</b>	VARIABLES1 : <code>collection(var-dvar)</code> VARIABLES2 : <code>collection(var-dvar)</code> SIZE_INTERVAL : <code>int</code>		
<b>Restrictions</b>	$ VARIABLES1  \geq  VARIABLES2 $ <code>required(VARIABLES1, var)</code> <code>required(VARIABLES2, var)</code> $SIZE\_INTERVAL > 0$		
<b>Purpose</b>	Let $N_i$ (respectively $M_i$ ) denote the number of variables of the collection <code>VARIABLES1</code> (respectively <code>VARIABLES2</code> ) that take a value in the interval $[SIZE\_INTERVAL \cdot i, SIZE\_INTERVAL \cdot i + SIZE\_INTERVAL - 1]$ . For all integer $i$ we have $M_i > 0 \Rightarrow N_i \geq M_i$ .		
<b>Example</b>	$(\langle 1, 9, 1, 8, 6, 2 \rangle, \langle 1, 0, 7, 7 \rangle, 3)$		
	<p>In the example, the third argument <code>SIZE_INTERVAL = 3</code> defines the following family of intervals <math>[3 \cdot k, 3 \cdot k + 2]</math>, where <math>k</math> is an integer. Consequently the values of the collection <code>VARIABLES2 = <math>\langle 1, 0, 7, 7 \rangle</math></code> are respectively located within intervals <math>[0, 2]</math>, <math>[0, 2]</math>, <math>[6, 8]</math>, <math>[6, 8]</math>. Therefore intervals <math>[0, 2]</math> and <math>[6, 8]</math> are respectively used 2 and 2 times. Similarly, the values of the collection <code>VARIABLES1 = <math>\langle 1, 9, 1, 8, 6, 2 \rangle</math></code> are respectively located within intervals <math>[0, 2]</math>, <math>[9, 11]</math>, <math>[0, 2]</math>, <math>[6, 8]</math>, <math>[6, 8]</math>, <math>[0, 2]</math>. Therefore intervals <math>[0, 2]</math>, <math>[6, 8]</math> and <math>[9, 11]</math> are respectively used 3, 2 and 1 times.</p> <p>Consequently, the <code>used_by_interval</code> constraint holds since, for each interval associated with the collection <code>VARIABLES2 = <math>\langle 1, 0, 7, 7 \rangle</math></code>, its number of occurrences within <code>VARIABLES1 = <math>\langle 1, 9, 1, 8, 6, 2 \rangle</math></code> is greater than or equal to its number of occurrences within <code>VARIABLES2</code>:</p> <ul style="list-style-type: none"> <li>• Interval <math>[0, 2]</math> occurs 3 times within <math>\langle 1, 9, 1, 8, 6, 2 \rangle</math> and 2 times within <math>\langle 1, 0, 7, 7 \rangle</math>.</li> <li>• Interval <math>[6, 8]</math> occurs 2 times within <math>\langle 1, 9, 1, 8, 6, 2 \rangle</math> and 2 times within <math>\langle 1, 0, 7, 7 \rangle</math>.</li> </ul>		
<b>Typical</b>	$ VARIABLES1  > 1$ <code>range(VARIABLES1.var) &gt; 1</code> $ VARIABLES2  > 1$ <code>range(VARIABLES2.var) &gt; 1</code> $SIZE\_INTERVAL > 1$ $SIZE\_INTERVAL < \text{range}(VARIABLES1.var)$ $SIZE\_INTERVAL < \text{range}(VARIABLES2.var)$		

**Symmetries**

- Items of VARIABLES1 are [permutable](#).
- Items of VARIABLES2 are [permutable](#).
- An occurrence of a value of VARIABLES1.var that belongs to the  $k$ -th interval, of size SIZE\_INTERVAL, can be [replaced](#) by any other value of the same interval.
- An occurrence of a value of VARIABLES2.var that belongs to the  $k$ -th interval, of size SIZE\_INTERVAL, can be [replaced](#) by any other value of the same interval.

**Arg. properties**

- [Contractible](#) wrt. VARIABLES2.
- [Extensible](#) wrt. VARIABLES1.
- [Aggregate](#): VARIABLES1(union), VARIABLES2(union), SIZE\_INTERVAL(id).

**Reformulation**

The `used_by_interval`( $\langle \text{var} - U_1 \text{ var} - U_2, \dots, \text{var} - U_{|\text{VARIABLES1}|} \rangle, \langle \text{var} - V_1 \text{ var} - V_2, \dots, \text{var} - V_{|\text{VARIABLES2}|} \rangle, \text{SIZE\_INTERVAL}$ ) constraint can be expressed by introducing  $|\text{VARIABLES1}| + |\text{VARIABLES2}|$  *quotient* variables

$$U_i = \text{SIZE\_INTERVAL} \cdot P_i + R_i, R_i \in [0, \text{SIZE\_INTERVAL} - 1] \quad (i \in [1, |\text{VARIABLES1}|]),$$

$$V_i = \text{SIZE\_INTERVAL} \cdot Q_i + S_i, S_i \in [0, \text{SIZE\_INTERVAL} - 1] \quad (i \in [1, |\text{VARIABLES2}|]),$$

in term of a conjunction of  $|\text{VARIABLES2}|$  reified constraints of the form:

$$\sum_{1 \leq j \leq |\text{VARIABLES1}|} (Q_i = P_j) \geq \sum_{1 \leq j \leq |\text{VARIABLES2}|} (Q_i = Q_j) \quad (i \in [1, |\text{VARIABLES2}|]).$$

**Used in**

[k\\_used\\_by\\_interval](#).

**See also**

[implied by: same\\_interval](#).

**soft variant:** [soft\\_used\\_by\\_interval\\_var](#) (*variable-based violation measure*).

**specialisation:** [used\\_by](#) (*variable/constant replaced by variable*).

**system of constraints:** [k\\_used\\_by\\_interval](#).

**Keywords**

**characteristic of a constraint:** [sort based reformulation](#).

**constraint arguments:** [constraint between two collections of variables](#).

**modelling:** [inclusion, interval](#).

<b>Arc input(s)</b>	VARIABLES1 VARIABLES2
<b>Arc generator</b>	$\text{PRODUCT} \mapsto \text{collection}(\text{variables1}, \text{variables2})$
<b>Arc arity</b>	2
<b>Arc constraint(s)</b>	$\text{variables1.var}/\text{SIZE\_INTERVAL} =$ $\text{variables2.var}/\text{SIZE\_INTERVAL}$
<b>Graph property(ies)</b>	<ul style="list-style-type: none"> <li>• for all connected components: <math>\text{NSOURCE} \geq \text{NSINK}</math></li> <li>• <math>\text{NSINK} =  \text{VARIABLES2} </math></li> </ul>

**Graph model**

Parts (A) and (B) of Figure 5.785 respectively show the initial and final graph associated with the **Example** slot. Since we use the **NSOURCE** and **NSINK** graph properties, the source and sink vertices of the final graph are stressed with a double circle. Since there is a constraint on each connected component of the final graph we also show the different connected components. Each of them corresponds to an equivalence class according to the arc constraint. Note that the vertex corresponding to the variable that takes value 9 was removed from the final graph since there is no arc for which the associated equivalence constraint holds. The `used_by_interval` constraint holds since:

- For each connected component of the final graph the number of sources is greater than or equal to the number of sinks.
- The number of sinks of the final graph is equal to  $|\text{VARIABLES2}|$ .

**Signature**

Since the initial graph contains only sources and sinks, and since sources of the initial graph cannot become sinks of the final graph, we have that the maximum number of sinks of the final graph is equal to  $|\text{VARIABLES2}|$ . Therefore we can rewrite  $\text{NSINK} = |\text{VARIABLES2}|$  to  $\text{NSINK} \geq |\text{VARIABLES2}|$  and simplify  $\overline{\text{NSINK}}$  to  $\text{NSINK}$ .

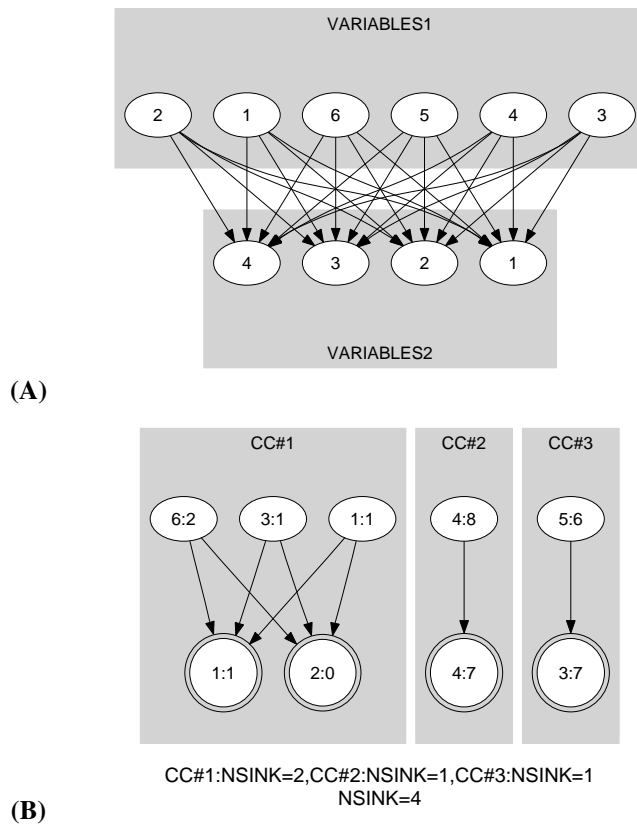


Figure 5.785: Initial and final graph of the used\_by\_interval constraint