

## 5.71 clique

	DESCRIPTION	LINKS	GRAPH
<b>Origin</b>	[159]		
<b>Constraint</b>	clique(SIZE_CLIQUÉ, NODES)		
<b>Arguments</b>	SIZE_CLIQUÉ : dvar NODES : collection(index-int, succ-svar)		
<b>Restrictions</b>	SIZE_CLIQUÉ $\geq 0$ SIZE_CLIQUÉ $\leq  \text{NODES} $ required(NODES, [index, succ]) NODES.index $\geq 1$ NODES.index $\leq  \text{NODES} $ distinct(NODES, index) NODES.succ $\geq 1$ NODES.succ $\leq  \text{NODES} $		
<b>Purpose</b>	Consider a digraph $G$ described by the NODES collection: to the $i^{\text{th}}$ item of the NODES collection corresponds the $i^{\text{th}}$ vertex of $G$ ; To each value $j$ of the $i^{\text{th}}$ succ variable corresponds an arc from the $i^{\text{th}}$ vertex to the $j^{\text{th}}$ vertex. Select a subset $\mathcal{S}$ of the vertices of $G$ that forms a clique of size SIZE_CLIQUÉ (i.e., there is an arc between each pair of distinct vertices of $\mathcal{S}$ ).		
<b>Example</b>	$\left( 3, \left\langle \begin{array}{ll} \text{index} - 1 & \text{succ} - \emptyset, \\ \text{index} - 2 & \text{succ} - \{3, 5\}, \\ \text{index} - 3 & \text{succ} - \{2, 5\}, \\ \text{index} - 4 & \text{succ} - \emptyset, \\ \text{index} - 5 & \text{succ} - \{2, 3\} \end{array} \right\rangle \right)$		
	The clique constraint holds since the NODES collection depicts a clique involving 3 vertices (namely vertices 2, 3 and 5) and since its first argument SIZE_CLIQUÉ is set to the number of vertices of this clique.		
<b>Typical</b>	SIZE_CLIQUÉ $\geq 2$ SIZE_CLIQUÉ $<  \text{NODES} $ $ \text{NODES}  > 2$		
<b>Symmetry</b>	Items of NODES are <a href="#">permutable</a> .		
<b>Arg. properties</b>	<a href="#">Functional dependency</a> : SIZE_CLIQUÉ determined by NODES.		
<b>Algorithm</b>	[159], [347, 348]. The algorithm for finding maximum cliques in an undirected graph of C. Bron and J. Kerbosch [88] was adapted by J.-C. Régin to the context of constraint programming in his papers.		

**See also**

**common keyword:** `link_set_to_booleans` (*constraint involving set variables, can be used for channelling*).

**used in graph description:** `in_set`.

**Keywords**

**constraint arguments:** constraint involving set variables.

**constraint type:** graph constraint.

**final graph structure:** symmetric.

**modelling:** functional dependency.

**problems:** maximum clique.

<b>Arc input(s)</b>	NODES
<b>Arc generator</b>	<code>CLIQUE(≠) ↦ collection(nodes1, nodes2)</code>
<b>Arc arity</b>	2
<b>Arc constraint(s)</b>	<code>in_set(nodes2.index, nodes1.succ)</code>
<b>Graph property(ies)</b>	<ul style="list-style-type: none"> <li>• <b>NARC</b> = SIZE_CLIQUÉ * SIZE_CLIQUÉ – SIZE_CLIQUÉ</li> <li>• <b>NVERTEX</b> = SIZE_CLIQUÉ</li> </ul>
<b>Graph class</b>	<b>SYMMETRIC</b>

**Graph model**

Note the use of *set variables* for modelling the fact that the vertices of the final graph have more than one successor: The successor variable associated with each vertex contains the successors of the corresponding vertex.

Part (A) of Figure 5.175 shows the initial graph from which we start. It is derived from the set associated with each vertex. Each set describes the potential values of the `succ` attribute of a given vertex. Part (B) of Figure 5.175 gives the final graph associated with the **Example** slot. Since we both use the **NARC** and **NVERTEX** graph properties, the arcs and the vertices of the final graph are stressed in bold. The final graph corresponds to a clique containing three vertices.

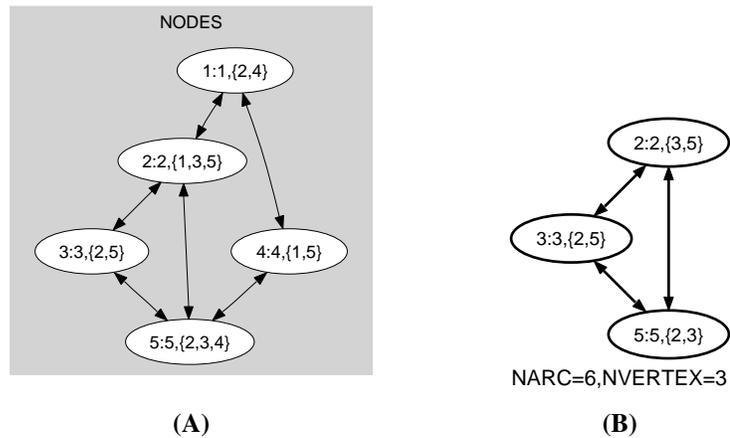


Figure 5.175: Initial and final graph of the `clique` set constraint

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